Performance evaluation of comprehensive bandwidth utilization for 10-gigabit passive optical network

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ABSTRACT
Bandwidth allocation during upstream transmission is crucial to determine the efficiency and performance of a XG-PON. For XG-PON, bandwidth assignment is done based on T-CONT which represents a traffic class as per ITU recommendation. DBA scheme used in this paper is based on CBU to assign bandwidth to ONUs based on the T-CONT supporting QoS as per SLA. In this paper, CATV traffic is used as traffic generator which used for generation of Ethernet frames and results showed expected trend of mean upstream delay for traffic class T2, T3 and T4 as compared to recommended value which is below 1.5ms. These results prove that CBU can also be implemented on real time traffic.

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1. INTRODUCTION
Information and communication technology (ICT) contribute immensely to social and economic improvements, such as higher employment and productivity, increasing access to a higher quality of life. The energy consumption from the expanding use of ICT is unsustainable with present drivers, and it will impact heavily on the future climate change. The expanding of ICT gives an enormous impact on electricity consumption as well as carbon dioxide (CO2) emissions. Over the last several years, a significant amount of research effort in improving energy saving performance in ICT equipment have been done. The fact that motivated research from industries and academia is the stupendously growing footprint of ICT in increasing energy demand. The electricity and heat production contribute 25% of global greenhouse gas (GHG) emissions reported by Intergovernmental Panel on Climate Change (IPCC) in 2014 [1]. In [2], it is reported that currently ICT is responsible between 5% to 9% of the global electricity consumption. Also, the growth rate is estimated to be 9% per year and it may rise up to 20% in year 2030.

Passive Optical Network (PON) is one of access network technologies that has been standardized by International Telecommunication Union (ITU-T) and Institute of Electrical and Electronics Engineers (IEEE). PON carries multiple services such as plain old telephone service (POTS), voice over IP (VoIP), data and video. PON consists of Optical Line Terminal (OLT) at the central office (CO), Optical Network Unit (ONU) and Optical Distribution Network (ODN) that is connecting the OLT and ONU using shared optical fiber and passive optical splitter. It utilizes the tree topology where OLT is the root of the tree. Passive elements in PON that do not require power-feeding makes it favorable as compared to other access networks to conserve energy and power consumption. The burst nature of data transmission between OLT–ONUs
contributes to significant amount of power consumption in PON. Generally, ONU's contribute major power consumption to the overall PON since OLT is connected to ONUs in point-to-multipoint connection. Data transmission between ONUs–OLT is Time Division Multiplexed (TDM), where every ONU is allocated a time slot to transmit data to OLT, avoiding collision.

As demands for higher bandwidth increases, the need for PON technologies to expand rises too. Specifically, Gigabit PON (GPON) has been extended to develop 10-gigabit capable PON (XG-PON) which supports data rate of 10Gbps in downstream (DS) and 2.5Gbps in upstream (US) and has a synchronous fixed frame duration of 125µs [3]. As mentioned earlier, US transmission of PON is in TDM, therefore bandwidth allocation algorithm is performed at the OLT to distribute bandwidth efficiently to all ONUs by allocating non-overlapping transmission slots to ONUs. Bandwidth allocation can be either static bandwidth allocation (SBA) or dynamic bandwidth allocation (DBA). However, bandwidth assignment in ITU PON is based on traffic container represented by T-CONT and DBA is normally practical to enrich the efficiency of a network. T-CONT 1 has a fixed bandwidth assignment. T-CONT 2 and T-CONT 3 have both assured and non-assured bandwidth assignment and T-CONT 4 has best-effort assignment of bandwidth. This paper presents an improvement of DBA scheduling using Comprehensive Bandwidth Utilization (CBU) [4] that is tested on Cable-TV (CATV) traffic [5] and results show that the US delay improved for T2, T3 and T4. In section 2, we will review DBA scheme for ITU PON. In section 3, we present our simulation setup. In section 4, we discussed about the results obtained from simulation. In section 5, we conclude our findings and proposed future work.

2. RELATED WORKS

The expansions of global access network has created more research opportunities especially in sustainable green communication area mapped to Quality of Service (QoS). Power consumption in optical access network is becoming a critical issue as demand for high speed broadband services increases. Power consumption in PON mainly caused by broadcast mechanism in downstream transmission where ONUs receive downstream frames from OLT. ONU’s need to be in active mode to receive DS frame and process it. Some of DS frames are irrelevant and will be discarded [6]. This scenario contributes to large power consumption in ONUs thus increasing total power consumption of a PON. Energy conservation techniques in PON can be looked into two distinct perspectives; physical layer based and MAC layer based. In [7] the OLT energy consumption has been studied due to the fact that the provisioning data rate keeps increasing which causes larger energy consumption. The main idea is to adapt the power-on OLT line cards in OLT chassis based on real time incoming traffic. Physical layer solutions are not favorable by network operators because they increase operation cost. MAC layer based solutions do not require additional equipment in the network and can be implemented on the existing network, thus it is more feasible and cost efficient. ITU has recommended power saving techniques in ONU for GPON/XG-PON that includes power shedding, dozing, deep sleep, cyclic sleep and watchful sleep [8].

Generally, performance measurement for DBA scheme is based on US delay because DBA schedules the time slot for ONUs to grant bandwidth based on traffic classes. The delay mainly comes from queuing time and idle time [9]. For energy efficient PON, minimum delays as an outcome of a good DBA is required as well. There are many DBA schemes presented on International Electrical and Electronics Engineers (IEEE) PON’s (e.g. EPON) [10–13], however cannot be used on ITU PONs because of the differences in layers. Additionally, for ITU PONs, OLT and ONUs are required to send frames every 125µs, while for IEEE PONs, OLT and ONUs only send frames when traffic arrived.

2.1. DBA schemes for ITU PON

Since OLT will grant bandwidth to ONUs, it needs to be aware of the queue reports for all ONUs. OLT maintains queue reports in the scenario of SR based DBA. ITU PON categorized traffic into four main types using traffic container (T-CONT) as shown in Table 1 [14]. Since T1 is a fixed bandwidth, it will not be considered in DBA. DBA is performed for allocation to T2, T3 and T4. The basis of DBA scheme for ITU PON started from [15] known as GigaPON Access Network (GIANT). There are two phases of bandwidth allocation namely Guaranteed Phase Allocation (GPA) and Surplus Phase Allocation (SPA). Four service parameters are calculated which are maximum and minimum service interval (Slmax and Slmin) and minimum and surplus allocation bytes (ABmin and ABsur). This Each queue has available byte (AB) counter and count down counter for faster bandwidth allocation. For fairness, the countdown counter follows round robin manner. The downside of it is that it must wait for count down counter to be expired, only will the available byte counter be recharged.
To overcome the problem, Immediate Allocation Colorless Grant (IACG) [16] allows GPA and SPA to be executed in every downstream frame and send bandwidth grants to ONUs. It improves US delay as compared to GIANT because bandwidth grants are sent every downstream frame but this causes bandwidth redundancy due to remaining bandwidth utilization. It improves US delay performance as compared to GIANT. In [17], the efficient bandwidth utilization (EBU) improves the utilization of remaining bandwidth by allowing the available byte counter to be negative. If the available byte counter is positive then it will adjust the unused bandwidth to be allocated among the same traffic classes, and if it is negative, it means that all traffic classes are overloaded. This adjustment leads to low bandwidth availability during a SI. Improving remaining bandwidth utilization is also presented in [18] where at the end of DBA cycle, remaining bandwidth are allocated to a fixed percentage of 40%, 35% and 25% for T2, T3 and T4 respectively. Another simple method presented in [19] uses a single available byte counter and single down counter for multiple queues and in [20], a repeated scheduling of assured bandwidth is introduced to ensure optimum bandwidth utilization. The work presented in [21] maximizes the bandwidth allocation to T2 during a SI even if it exceeds the provisioned AB. The results show an improvement in mean US delay for T2.

All of the above-mentioned DBA schemes use the four service parameters in their algorithms. According to [22], real data queuing information and independent from SI improves average delay by 40% for T2 as well as lowers bandwidth waste as compared to SI adaptive schemes. DBA method in [23] grouped four cycles to identify DBA computing time. Traffic prediction queue reports are presented in [24-26] demonstrated the NSR in GPON. In [18], [27] polling cycle time is improved and thus the idle time is also improved. A request-based polling time presented in [28] improves queuing delay such that every traffic class has its own polling order. The polling order is based on highest request bandwidth.

3. SIMULATION SETUP

A simulation model was developed in OMNET++ to evaluate the performance of CBU scheme. A single OLT is connected to 16 ONUs over a single optical fiber using a passive optical splitter. The DS traffic load is fixed to 0.1 and US traffic load varies from 0.1 to 1. Since T1 is allocated for fixed bandwidth, T1 is not considered in this paper. The simulation setup shown in Figure 1 shows a snapshot of our simulation testbed. All the important simulation parameters are tabulated in Table 2.

<table>
<thead>
<tr>
<th>TCONT</th>
<th>Type of Bandwidth</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Fixed bandwidth</td>
<td>VoIP, live video streaming</td>
</tr>
<tr>
<td>2</td>
<td>Assured bandwidth</td>
<td>Basic video/voice</td>
</tr>
<tr>
<td>3</td>
<td>Non-assured</td>
<td>VoIP/Video conferencing</td>
</tr>
<tr>
<td>4</td>
<td>Best effort</td>
<td>Instant messaging/emails</td>
</tr>
</tbody>
</table>
Table 2. Simulation parameters

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Values/Details</th>
</tr>
</thead>
<tbody>
<tr>
<td>US/DS Line rates</td>
<td>10Gbps/2.5Gbps</td>
</tr>
<tr>
<td>DS Traffic Load</td>
<td>Fixed to 0.01</td>
</tr>
<tr>
<td>US Traffic Load</td>
<td>Varied from 0.1 to 1</td>
</tr>
<tr>
<td>Bandwidth Assignment for T2</td>
<td>$AB_{\text{min}}=7812$ bytes with $SI_{\text{max}}=5$ (100 Mbps)</td>
</tr>
<tr>
<td>Bandwidth Assignment for T3</td>
<td>$AB_{\text{min}}=AB_{\text{idle}}=7812$ with $SI_{\text{max}}=SI_{\text{min}}=10$ (50 Mbps Assured and Non-assured portions).</td>
</tr>
<tr>
<td>Bandwidth Assignment for T4</td>
<td>$AB_{\text{min}}=15,624$ and $SI_{\text{max}}=10$ (100 Mbps)</td>
</tr>
</tbody>
</table>

4. RESULTS AND ANALYSIS

To measure the performance of a DBA scheme, mean US delay is taken into consideration, where it is a summation of $T_{\text{queueing}}$ and $T_{\text{idle}}$. $T_{\text{idle}}$ is waiting time of the frame until bandwidth is obtained and $T_{\text{queueing}}$ is waiting time of the frame in the queue after being reported to OLT [29]. The average US delay results for T2, T3 and T4 traffic are illustrated in Figure 2, 3, and 4. ITU-T recommendation in [8] recommends that XGPON system must accommodate services that require a maximum mean delay of 1.5ms. We compare the result with the ITU-T recommendation and observe a significant less delay for T2, T3 and T4 for every network load. The results are analyzed based on CATV traffic which used for generation of Ethernet frames and they show expected trends as compared to recommended mean delay.

T2 traffic class shows the least delay as network load increases due to its highest priority. This proves that the algorithm satisfies ITU standard which T2 and T3 are supposed to be given priority bandwidth allocation and T4 has the least priority. T4 traffic class has the highest US delay as network load increases. The increased of delay for T4 starts at network load 0.8 where it is already beyond XG-PON requirement. This is expected because T4 is best effort type of bandwidth and given the lowest priority with 25% of bandwidth assignment. T2 and T3 have the average US delay less than ITU-T recommendation.

Bandwidth wasted in CATV traffic shows a decreasing trend as network load increases as shown in Figure 5. The trend is acceptable because as network load increases, bandwidth demand increases hence there is not much bandwidth wasted per cycle. The results show that CBU can also be implemented on a real time traffic.

![Figure 2. Mean US delay for T2](image1)

![Figure 3. Mean US delay for T3](image2)
5. CONCLUSION

In this paper we have reviewed DBA schemes for ITU-PON also we have studied the DBA scheme for XG-PON using CBU to analyze the mean US delay for T2, T3 and T4. T1 is not included because it is a fixed bandwidth. Broadcom CATV is used as traffic generator to imitate real time traffic. The results summarized that CBU can be implemented to real time traffic as the US delay is below target value as recommended by ITU. The analysis on CBU DBA scheme can be further studied on uniform and triangular distribution, thus comparing the three types of distributions. This will be the next step of the research. In future, we will further improve CBU in terms of fairness as well as combining sleep assistive ONU to study the energy efficiency of XG-PON.

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